

## A Discrete Intermediate Value Theorem

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Many theories in mathematics (for instance, difference and differential equations) come in discrete and continuous versions. Indeed, an entire book, *Excursions in Calculus: An Interplay of the Continuous and the Discrete* by Robert M. Young (MAA, Washington, DC, 1992), has been devoted to this topic.

In this note, I give a discrete intermediate value theorem and apply it in an appealing induction proof.

The continuous intermediate value theorem is well known:

*If  $f$  is continuous on  $[a, b]$  and  $f(a)f(b) < 0$ , then  $f(x) = 0$  for some  $x$  in  $(a, b)$ .*

A discrete intermediate value theorem (DIVT) is:

*Let  $f$  be an integer-valued function defined on the integers in  $[m, n]$ . Suppose (as the equivalent of a continuity assumption) that  $|f(i) - f(i + 1)| \leq 1$  for  $m \leq i < n$ . If  $f(m)f(n) < 0$ , then  $f(x) = 0$  for some integer  $x$  in  $(m, n)$ .*

As an application of the DIVT, consider the set  $L$  of all strings of the symbols  $a$  and  $b$  that can be constructed from the null string by application of a finite sequence of the rules:

1. if  $\alpha \in L$  then  $a\alpha b \in L$  and  $b\alpha a \in L$ ;
2. if  $\alpha \in L$  and  $\beta \in L$  then  $\alpha\beta \in L$ .

It is clear from this definition that every element of  $L$  has equal numbers of  $a$ 's and  $b$ 's. But what about the converse?

Let  $\alpha$  be a string of length  $n$  with equal numbers of  $a$ 's and  $b$ 's. We use induction on  $n$  to show that  $\alpha$  is in  $L$ . The base case,  $n = 0$ , is the null string, which by definition is in  $L$ . For the inductive step, assume that  $n > 0$  and that any string of length less than  $n$  with equal numbers of  $a$ 's and  $b$ 's is in  $L$ .

First suppose that  $\alpha$  starts with  $a$  and ends with  $b$ . That is,  $\alpha = a\beta b$  for some string  $\beta$ . Then  $\beta$  has equal numbers of  $a$ 's and  $b$ 's. So, by the inductive assumption,  $\beta$  is in  $L$ . It now follows that  $\alpha$  is in  $L$  by rule 1. By the same argument, if  $\alpha$  starts with  $b$  and ends with  $a$ ,  $\alpha$  is in  $L$ .

Now suppose that  $\alpha$  starts and ends with  $a$  (and still has equal numbers of  $a$ 's and  $b$ 's). Here is where the DIVT can be used. Let

$$f(i) = \text{number of } a\text{'s} - \text{number of } b\text{'s}$$

within the first  $i$  symbols of  $\alpha$ . Then  $f(1) = 1$  and  $f(n) = -1$ .

Since  $f(i + 1)$  and  $f(i)$  always differ by 1,  $f$  satisfies the "continuity" assumption, and the DIVT says that  $f(i) = 0$  for some  $i$  between 1 and  $n - 1$ . It follows that  $\alpha = \beta\gamma$ , where  $\beta$  and  $\gamma$  are non-null strings having equal numbers of  $a$ 's and  $b$ 's. By the inductive assumption,  $\beta$  and  $\gamma$  are in  $L$  and, hence,  $\beta\gamma$  is in  $L$  by rule 2. Similarly, if  $\alpha$  starts and ends with  $b$ ,  $\alpha$  is in  $L$ .

